



Ex 3:

$$A = \begin{bmatrix} x & x & x & x \\ x & x & & \\ x & & x & \\ x & & & x \end{bmatrix} \quad \left\{ \begin{array}{l} \text{Arrowhead} \\ \beta_{\max} = 3 \\ \varepsilon = 6 \end{array} \right.$$

Instead of solving  $A\underline{x} = \underline{b}$ ; solve  $\underbrace{PAP^T}_P \underline{x} = \underline{Pb}$

Use for  $\alpha$ .

$$P = \begin{bmatrix} & & & \\ & & & \\ & & & \\ 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix}; \text{ Then } PAP = \begin{bmatrix} x & & & \\ & x & & \\ & & x & \\ & & & x \\ x & x & x & x \end{bmatrix} \quad \left\{ \begin{array}{l} \beta_{\max} = 3 \\ \varepsilon = 3 \end{array} \right.$$

Optimal ?  $\uparrow$

We will see:

- This was not optimal
- We need a systematic process to find P.

Some illustrative examples:

Ex 4:

$$\begin{bmatrix} x & x \\ x & x & x \\ & x & x & x \\ & & x & x & x \\ & & & x & x \end{bmatrix}$$

Draw graph, much as for extended Gershgorin's theorem, (to check irreducibility), but now lines only; no arrows.



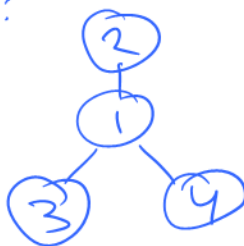
We can from graph read off

$$\left\{ \begin{array}{l} \beta_1 = \max(0, 1-2) = 0 \\ \beta_2 = \max(0, 2-1, 2-3) = 1 \\ \vdots \\ \beta_5 = \max(0, 5-4) = 1 \end{array} \right.$$

$$\Rightarrow \left\{ \begin{array}{l} \beta_{\max} = 1 \\ \varepsilon = 4 \end{array} \right.$$

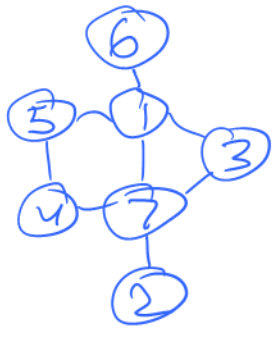
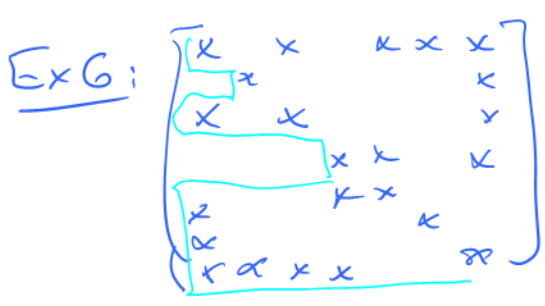
Ex 5: Arrowhead:

$$\begin{bmatrix} x & x & x & x \\ x & x & & \\ x & & x & \\ x & & & x \end{bmatrix}$$



$$\left\{ \begin{array}{l} \beta_1 = \max(0, 1-2, 1-3, 1-4) = 0 \\ \beta_2 = \max(0, 2-1) = 1 \\ \beta_3 = \max(0, 3-1) = 2 \\ \beta_4 = \max(0, 4-1) = 3 \end{array} \right.$$

$$\text{So (as before)} \left\{ \begin{array}{l} \beta_{\max} = 3 \\ \varepsilon = 6 \end{array} \right.$$



$$\beta_1 = 0$$

$$\beta_2 = 0$$

$$\beta_3 = 2$$

$$\beta_4 = 0$$

$$\beta_5 = 4$$

$$\beta_6 = 5$$

$$\beta_7 = 6$$

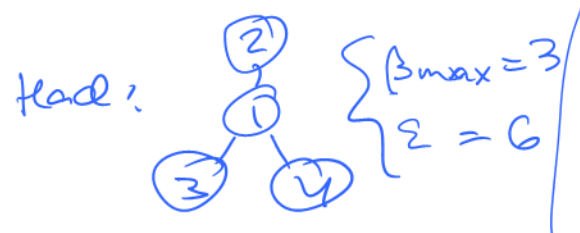
$$\left\{ \begin{array}{l} \beta_{max} = 6 \\ \Sigma = 17 \end{array} \right.$$

Note: Different  $\beta$  correspond to different orderings of the numbers 1, 2, ... within the graph.

Cuthill-McKee idea:

- Start 'far out' in the graph (max distance to other elements)
- Fill in 1, 2, 3 ... ensuring that neighbors to low elements get filled first.
- Reverse.

Ex 5 (re-visited) Arrowhead:



$$\left\{ \begin{array}{l} \beta_{max} = 3 \\ \Sigma = 6 \end{array} \right.$$



$$\beta_1 = \max(0, 1-2) = 0$$

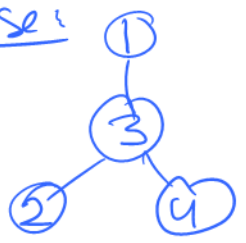
$$\beta_2 = \max(0, 2-1, 2-3, 2-4) = 1$$

$$\beta_3 = \max(0, 3-2) = 1$$

$$\beta_4 = \max(0, 4-2) = 2$$

$$\left\{ \begin{array}{l} \beta_{max} = 2 \\ \Sigma = 4 \end{array} \right.$$

Reverse:



$$\beta_1 = \max(0, 1-3) = 0$$

$$\beta_2 = \max(0, 2-3) = 0$$

$$\beta_3 = \max(0, 3-1, 3-2, 3-4) = 2$$

$$\beta_4 = \max(0, 4-3) = 1$$

$$\left\{ \begin{array}{l} \beta_{max} = 2 \\ \Sigma = 3 \end{array} \right.$$

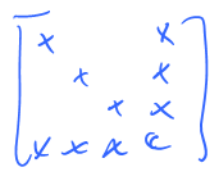
Sequence:

original



$$\left\{ \begin{array}{l} \beta_{max} = 3 \\ \Sigma = 6 \end{array} \right.$$

looked good



$$\left\{ \begin{array}{l} \beta_{max} = 3 \\ \Sigma = 3 \end{array} \right.$$

Cuthill-McKee



$$\left\{ \begin{array}{l} \beta_{max} = 2 \\ \Sigma = 4 \end{array} \right.$$

Reverse Cuthill-McKee



$$\left\{ \begin{array}{l} \beta_{max} = 2 \\ \Sigma = 3 \end{array} \right.$$

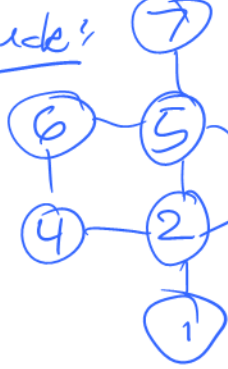
Theorem: In the reverse step:

→  $\beta_{max}$  always same

→  $\epsilon$  may decrease (never increase)

Ex 6: (re-visited)

C-Mck:

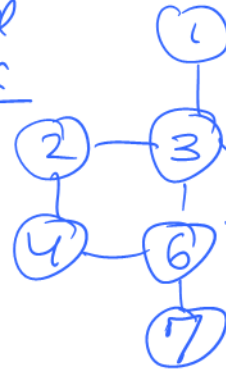


$\beta_1=0$   
 $\beta_2=1$   
 $\beta_3=1$   
 $\beta_4=2$   
 $\beta_5=3$   
 $\beta_6=2$   
 $\beta_7=2$

$\beta_{max}=3$   
 $\epsilon=11$

Reversed

C-Mck



$\beta_1=0$   
 $\beta_2=0$   
 $\beta_3=2$   
 $\beta_4=2$   
 $\beta_5=2$   
 $\beta_6=3$   
 $\beta_7=1$

$\beta_{max}=3$   
 $\epsilon=10$

Original  $\left\{ \begin{array}{l} \beta_{max}=6 \\ \epsilon=17 \end{array} \right.$

Typical: Larger scale problems  $\Rightarrow$  bigger gain.

Issues:

1. Reverse Cuthill-McKee very good, but not theoretically optimal.

2. Discussion here limited to A symmetric.

For non-symmetric, might need to pivot w/ conditioning as well.

3. Cost increase for larger graphs.  $\left\{ \begin{array}{l} \text{Open to a big issue!} \\ \text{symmnc very fast.} \end{array} \right.$