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# **RBF** Partition of Unity

February 28, 2018

### **Outline**

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### **Problem**

- Solving partial differential equations (PDEs) or interpolating data with global radial basis functions (RBFs) can lead to large node sets
- Therefore the dependence of computational time as a function of the number of nodes (N) becomes important
- Using N nodes with Global RBF  $\phi(|\mathbf{x}_i \mathbf{x}_i|)$ , A is a dense N by N matrix

$$A = \begin{bmatrix} \phi(|\mathbf{x}_0 - \mathbf{x}_0|) & \phi(|\mathbf{x}_0 - \mathbf{x}_1|) & \dots & \phi(|\mathbf{x}_0 - \mathbf{x}_{N-1}|) \\ \phi(|\mathbf{x}_1 - \mathbf{x}_0|) & \phi(|\mathbf{x}_1 - \mathbf{x}_1|) & \dots & \phi(|\mathbf{x}_0 - \mathbf{x}_{N-1}|) \\ \vdots & \vdots & \ddots & \vdots \\ \phi(|\mathbf{x}_{N-1} - \mathbf{x}_0|) & \phi(|\mathbf{x}_{N-1} - \mathbf{x}_1|) & \dots & \phi(|\mathbf{x}_{N-1} - \mathbf{x}_{N-1}|) \end{bmatrix}$$

Conclusion

### **Problem**

- For a dense N by N matrix A
- Getting weights requires solving a system of linear equation  $\mathbf{A}\lambda = \mathbf{u}$ 
  - O(N³) operation
- Applying the Differentiation Matrix (DM) is a matrix multiplication
  - O(N²) operation
- The resulting methods are computational difficult for large node sets
- RBF Partition of Unity (RBF-PUM) methods look to avoid these issues
- An additional leading approach for attacking these issues is RBF-FD

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### Brief RBF-PUM history

- The idea of RBF-PUM was original proposed by Babuška and Melenk (1997)
- Used for interpolation with compactly supported RBFs by Wedland (2002)
- Further discussion in Fasshauer's book (2007)
- Explored for interpolating non uniform data sets in Cavoretto, De Rossi et. al. (2012-Present)
- Used to solve PDEs in E. Larsson et. al. (2012-present) and J. Ahlkrona and V. Shcherbakov (2017-present)

Conclusions

# Partition of Unity

- Rather than taking N global nodes, the space is partitioned
- The partitions overlap, however the sum of their weights are unity throughout the whole space.
- The resulting matrix becomes sparse
- Aλ = u solution and matrix multiplication can be done significantly faster for sparse matrices.

Conclusions

# Partition of Unity

Partition of Unity starts with a set of patches  $\Omega_j$ , j = 1, ..., P that cover a domain  $\Omega$ 

$$\bigcup_{j=1}^P \Omega_j \supseteq \Omega$$

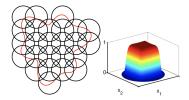


Fig. 3.1. To the left, the red curve is the outline of the domain  $\Omega_S$  and the black circles are the boundaries of the overlapping circular patches  $\Omega_j$ ,  $j = 1, \dots, P$ . To the right, a partition of unity weight function  $w_j$  for one of the interior patches is shown.

# Partition of Unity

Each patch has a weight  $\omega_i$  such that

$$\sum_{j=1}^{P} \omega_j(\mathbf{x}) = 1, \forall \mathbf{x} \in \Omega$$

$$\omega_j(\mathbf{x}) = \frac{\phi_j(\mathbf{x})}{\sum_{i=1}^P \phi(\mathbf{x})}$$

One set of generating functions  $\phi(r)$  (Wendland functions)

$$\phi_j(r) = (4r+1)(1-r^4)_+$$

pertain to a patch with center  $\mathbf{c}_i$  and radius  $\rho_i$ 

$$\phi_j(\mathbf{x}) = \phi_j \left( \frac{|\mathbf{x} - \mathbf{c}_j|}{\rho_j} \right)$$

giving an interpolate from local interpolates  $u_i(\mathbf{x})$ 

$$u(\mathbf{x}) = \sum_{j=1}^{P} \omega_j(\mathbf{x}) u_j(\mathbf{x})$$

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### RBF-PUM C-RBF-PUM LS-RBF-PUM

Conclusion

To move from PUM to RBF-PUM one introduces a RBF  $\Phi(r)$  such as a Gaussian

$$\Phi(r) = e^{-(\epsilon r)^2}$$

and a set of node points  $\mathbf{x}_i$  where  $u(\mathbf{x}_i)$  is known. Using the global RBF tools one gets the local interpolates  $u_j(\mathbf{x})$  such that

$$u_j(\mathbf{x}) = \sum_{i=1}^{n_j} \lambda_i^j \Phi(|\mathbf{x} - \mathbf{x}_i|)$$

with the  $\lambda_i^j$  obtained using Global RBF in the particular patch. (RBF-QR)

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### RBF-PUM

LS-RBF-PUM Patch Sets

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The same framework can be extended to a differential operator  $\boldsymbol{\mathcal{L}}$  such that

$$\mathcal{L}u(\mathbf{x}) = \sum_{j=1}^{P} \mathcal{L}(\omega_{j}(\mathbf{x})u_{j}(\mathbf{x}))$$

This allows for the use of RBF-PUM to solve PDEs

### **C-RBF-PUM**

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Global node set

$$X = \{\boldsymbol{x}_k\}_{k=1}^N$$

Consisting of interior nodes

$$X^i = \{ \mathbf{x}_k \in X : \mathbf{x}_k \in \Omega \}$$

and boundary nodes

$$X^b = \{ \mathbf{x}_k \in X : \mathbf{x}_k \in \partial \Omega \}$$

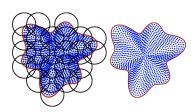


Fig. 3.2. A global node set and patches in  $\Omega_S$  (left), and evaluation points for collocation (same as node points) (right) for C-RBF-PUM.

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C-RBF-PUM LS-RBF-PUM Patch Sets

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### LS-RBF-PUM

- Start with identical node sets  $X_i$  with respect to a patch  $\Omega_i$
- · Global set gives

$$X = \bigcup_{j=1}^{P} X_{j}$$

- The uniformity of  $X_i$  lead to only one linear system solve
- Decouples evaluation points  $Y^i \subset Y$  and  $Y^b \subset Y$

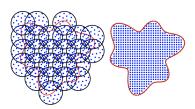
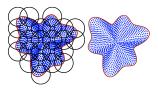


Fig. 3.3. Patches with identically distributed local node sets covering the domain  $\Omega_S$  (left), and loast squares evaluation points on a Cartesian grid in the interior and uniform with respect to are length on the boundary (right) for LS-RBF-PUM.

LS-RBF-PUM

# Comparison



as node points) (right) for C-RBF-PUM.

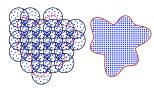


Fig. 3.3. Patches with identically distributed local node sets covering the domain Ω<sub>2</sub> (left), and FIG. 3.2. A global node set and patches in \Omega\_S (left), and evaluation points for collocation (same least squares evaluation points on a Cartesian grid in the interior and uniform with respect to arc length on the boundary (right) for LS-RBF-PUM.

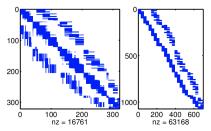


Fig. 3.4. The structure of the matrix L for a problem defined over  $\Omega_S$  with P = 24 patches, box size H=0.6, and overlap  $\delta=0.2$ . For C-RBF-PUM (left),  $13 \le n_j \le 42$ ,  $h\approx 0.12$ , and N=321, and for LS-RBF-PUM (right), n = 28,  $h \approx 0.14$ , N = 700, M = 1073, and  $\beta = M/N \approx 1.5$ .

C-RBF-PUM LS-RBF-PUM Patch Sets

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## Comparison

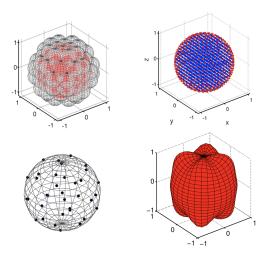


FIG. 5.10. Patches on the spherical domain  $\Omega_U$  (top left), quasi uniformly distributed least squares points inside  $\Omega_U$  and on the surface  $\partial\Omega_U$  (top right), a single patch (enlarged) with n=35local node points (bottom left), and the star shaped domain  $\Omega_Q$  (bottom right).

DUM

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### Patch Sets

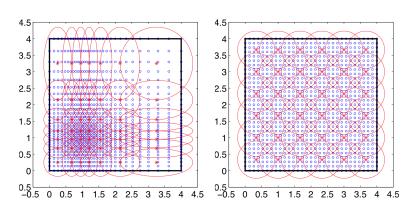


Fig. 12. Discretization of the square domain with oval and circle patches

A. Safdari-Vaighani, A.R.H. Heryudono, and E. Larsson (2015)

Conclusions

### Conclusions

- RBF-PUM was proposed in 1997
- C-RBF-PUM and LS-RBF-PUM algorithms were presented
- Expended from 2D to 3D with spheres
- Non uniform patch sets have been shown
- The algorithms presented are by no means exhaustive

### Questions

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Thank you for your attention!