

## Lecture #21

**1** We now consider the trigonometric interpolation problem. We seek for a function of the following form

$$p(x) = a_0 + \sum_{j=1}^n a_j \cos(jx) + b_j \sin(jx).$$

If  $a_n$  and  $b_n$  are not both zero, then this is called a trigonometric polynomial of degree  $n$ .

We've implicitly given a basis for our function space (dimension  $2n + 1$ )

$$\{1, \cos(x), \sin(x), \dots, \cos(nx), \sin(nx)\}.$$

We should really first consider whether these functions are linearly independent; otherwise the interpolation problem posed here might have multiple solutions (if it has any). I will assume you already know that these functions are a basis for their span.

We will reduce the trig interpolation problem to the polynomial one. Note that

$$e^{ijx} = \cos(jx) + i \sin(jx)$$

$$\cos(jx) = \frac{e^{ijx} + e^{-ijx}}{2}, \quad \sin(jx) = \frac{e^{ijx} - e^{-ijx}}{2i}$$

so

$$p(x) = \sum_{j=-n}^n \tilde{c}_j e^{ijx} = \sum_{j=-n}^n \tilde{c}_j (e^{ix})^j$$

$$\tilde{c}_0 = a_0, \quad \tilde{c}_j = \frac{1}{2}(a_j + ib_j), \quad \tilde{c}_{-j} = \tilde{c}_j^* \equiv a_j = \tilde{c}_j + \tilde{c}_{-j}, \quad b_j = \frac{\tilde{c}_j - \tilde{c}_{-j}}{i}$$

$$\tilde{p}(z) = \sum_{j=-n}^n \tilde{c}_j z^j, \quad p(x) = \tilde{p}(z) \text{ when } z = e^{ix}.$$

$$P(z) = z^n \tilde{p}(z) = \sum_{j=0}^{2n} c_j z^j, \quad \tilde{c}_{j-n} = c_j$$

$P$  is a polynomial of degree  $\leq 2n$ . Because there is a one-to-one relationship between the  $c_j$  and the  $a_j, b_j$ , the trig interpolation problem is equivalent a (complex) polynomial interpolation problem. The latter has a unique solution whenever the  $z_j$  are distinct; what does this mean for the  $x_j$ ? Since  $z = e^{ix}$ , the  $z_j$  will be distinct whenever the  $x_j$  are distinct modulo  $2\pi$ . Typically we would first rescale all of our  $x_j$  so that they lie in the interval  $[0, 2\pi)$ .

**2** The above shows that the trig interpolation problem has a unique solution provided that the  $x_j$  are distinct modulo  $2\pi$ ; what about computing the solution? For non-equispaced nodes you can just solve the associated polynomial problem using Newton Divided Differences, and then re-map the coefficients  $c_j$  to the  $a_j, b_j$ . If the nodes are equispaced there are better methods.

WLOG, let (not the same as Atkinson)

$$[0, 2\pi) \ni x_j = j \frac{2\pi}{2n+1}, \quad z_j = e^{ix_j}, \quad j = 0, \dots, 2n.$$

Consider the Vandermonde matrix of the associated polynomial interpolation problem.

$$P(z_j) = \sum_{k=0}^{2n} c_k z_j^k = c_0 + c_1 z_j + \dots + c_{2n} z_j^{2n} = y_j$$

$$\begin{bmatrix} 1 & z_0 & \dots & z_0^{2n} \\ 1 & z_1 & \dots & z_1^{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & z_{2n} & \dots & z_{2n}^{2n} \end{bmatrix}.$$

We will now show that the columns of this matrix are orthogonal with respect to the standard complex dot product. A single column has the form

$$\mathbf{v}_k = \mathbf{v}^{(k)}; \quad \mathbf{v}_j^{(k)} = z_j^k, \quad j = 0, \dots, 2n, \quad k = 0, \dots, 2n.$$

The complex dot product of two columns is

$$\mathbf{v}^{(k)} \cdot \mathbf{v}^{(l)} = \sum_{j=0}^{2n} z_j^k \overline{z_j^l}$$

Now use

$$z_j^k = \exp\left\{jk \frac{2\pi i}{2n+1}\right\}$$

$$\mathbf{v}^{(k)} \cdot \mathbf{v}^{(l)} = \sum_j \exp\left\{j \frac{2\pi i}{2n+1} (k-l)\right\} = \sum_{j=0}^{2n} \omega(k-l)^j, \quad \omega(k-l) = \exp\left\{(k-l) \frac{2\pi i}{2n+1}\right\}$$

We can use the geometric series formula

$$\sum_{j=0}^{2n} \omega(k-l)^j = \frac{1 - \omega(k-l)^{2n+1}}{1 - \omega(k-l)}.$$

Now consider

$$\omega(k-l)^{2n+1} = \exp\left\{(k-l) \frac{2\pi(2n+1)i}{2n+1}\right\} = 1.$$

If  $k-l=0$  we just have

$$\mathbf{v}^{(k)} \cdot \mathbf{v}^{(k)} = 2n+1.$$

The Vandermonde system can be solved easily

$$\mathbf{V}\mathbf{c} = \mathbf{y}, \quad \mathbf{c} = \frac{\mathbf{V}^* \mathbf{y}}{2n+1}.$$

The matrix  $\mathbf{V}^*$ , up to normalization, is the ‘discrete Fourier transform.’ There are fast algorithms to apply it to vectors (i.e. costing less than  $\mathcal{O}(n^2)$  flops) called the Fast Fourier Transform.

**3 Accuracy/convergence.** Theorem 3.6 in Atkinson. If  $f$  is  $2\pi$  periodic and continuous then the error in the trig interpolation (with equispaced nodes) is bounded by

$$\|f - p_n\|_\infty \leq c \ln(n+2) \rho_n(f)$$

where  $\rho_n(f)$  is the error in the best-possible approximation (in the  $\infty$  norm) of  $f$  using a trig polynomial of degree  $\leq n$ . The constant  $c$  is independent of  $f$  and  $n$ . Of course this doesn’t imply convergence unless you know how  $\rho_n$  scales with  $n$ . We’ll come back to equispaced trig interpolation from a different perspective later, and return to the question of errors & convergence.